# Department of Computer Science University of Cyprus



### **EPL646 – Advanced Topics in Databases**

## Lecture 16

Big Data Management VI (MapReduce Programming)

Credits: Pietro Michiardi (Eurecom): Scalable Algorithm Design,
Apache MapReduce Tutorial

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http://www.cs.ucy.ac.cy/~dzeina/courses/epl646

## Outline of Lecture



- MapReduce "Hello World" Program Explained
  - Wordcount in MR, Example Execution, Pseudocode
  - Mean Computation in MR, JAVA API Preview
- Operational Issues:
  - What to configure and what not
- Combiners and In-Memory Combiners
- Relational Operators in MR
  - Selection/Projection
  - Union / Intersection / Set Difference
  - Join /Aggregation

# Introduction to Hadoop Programming



- In the previous lecture we learnt how a MapReduce program executes in a Hadoop Environment without actually seeing the program.
- In this lecture we will learn more about the basic principles on how to write MapReduce Programs in Hadoop.
- To validate some of the ideas in this lecture, ensure that Hadoop is installed, configured and running. More details:
  - Single Node Setup for first-time users.
  - Cluster Setup for large, distributed clusters.
- In our laboratory we will use a Single Node Setup (consult the image that has been circulated by the TA).
  - Hadoop v2 requires Java 7 or greater
     — our labs & assignments ☺.
  - Hadoop v3 (Jan. 17, alpha) requires Java 8 will not be seen further ☺
    - New Features: HDFS Erasure encoding, YARN v2 Timeline service (HBase store) Opportunistic containers, 2 Namenodes, default ports changed, Filesystem Connectors (e.g., Microsoft Azure Data Lake), Intra-DataNode Balancer

# MapReduce "Hello World" (WordCount 1/2)



```
import java.io.IOException;
import java.util.StringTokenizer:
                                                                 As of JAVA 7 Generic Example: no way to verify.
                                                                 at compile time, how the class is used (e.g., as
import org.apache.hadoop.conf.Configuration:
                                                                 import org.apache.hadoop.fs.Path:
                                                                  public class Box {
import org.apache.hadoop.io.IntWritable;
                                                                      private Object object:
import org.apache.hadoop.io.Text;
                                                                      public void set(Object object) { this.object = object; }
import org.apache.hadoop.mapreduce.Job:
                                                                     public Object get() { return object; }
import org.apache.hadoop.mapreduce.Mapper;
import org.apache.hadoop.mapreduce.Reducer:
import org.apache.hadoop.mapreduce.lib.input.FileInputFormat:
import org.apache.hadoop.mapreduce.lib.output.FileOutputFormat:
public class WordCount {
 public static class TokenizerMapper extends Mapper<Object, Text, Text, IntWritable> {
     private final static IntWritable one = new IntWritable(1); // optimized serialization of JAVA.Integer() class
     private Text word = new Text();
                                   Input(k,v)
     public void map(Object key, Text value, Context context) throws IOException, InterruptedException {
        StringTokenizer itr = new StringTokenizer(value.toString());
                                                                               Applications may override
        while (itr.hasMoreTokens()) {
                                                                               the run(org.apache.hadoop.mapreduce.Mapp
          word.set(itr.nextToken()); context.write(word, one);
                                                                               er.Context) method to exert greater control
                                                                              on map processing e.g. field delimiters, etc.
     } // map
                                             Output(k',v')
```

# MapReduce "Hello World" (WordCount 2/2)



```
Input(k.v)
public static class IntSumReducer extends Reducer<Text.IntWritable.Text.IntWritable> {
 private IntWritable result = new IntWritable();
 public void reduce(Text key, Iterable<IntWritable> values, Context context) throws IOException. InterruptedException {
  int sum = 0:
  for (IntWritable val : values) {
    sum += val.get():
  result.set(sum):
  context.write(key, result);
             Output(k',v')
public static void main(String[] args) throws Exception {
 Configuration conf = new Configuration();
 Job job = Job.getInstance(conf, "word count");
 job.setJarByClass(WordCount.class);
 job.setMapperClass(TokenizerMapper.class);
 job.setCombinerClass(IntSumReducer.class);
 job.setReducerClass(IntSumReducer.class);
 job.setOutputKeyClass(Text.class);
 job.setOutputValueClass(IntWritable.class);
 FileInputFormat.addInputPath(job, new Path(args[0]));
 FileOutputFormat.setOutputPath(job, new Path(args[1]));
 System.exit(job.waitForCompletion(true) ? 0 : 1);
```

> Implementing this interface allows an object to be the target of the "for-each loop" statement. (as of 1.5)

Output(k'.v')

As of JAVA 5 Enhanced For Loop: iterate through all the elements of a *Collection*.

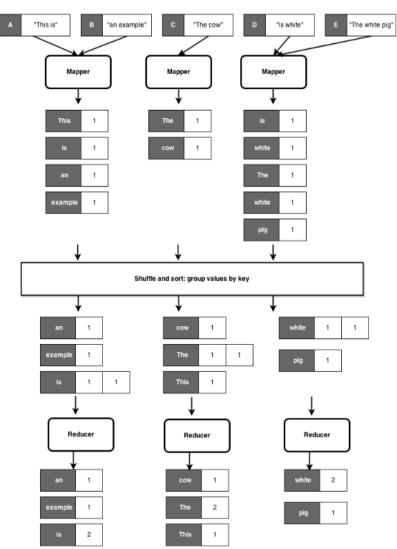
> Cleanup(), setup() are not mandatory in Mapper (see next slide)!

Multi-threading is possible with MultithreadedMapper when the Mapper is **not CPU bound** (the time to complete the task is **not** determined by the slow CPU but the Mapper logic). => If using more CPU power would speedup the program execution.

To-5

# Execution (Wordcount, Anagram)





In laboratory you saw the anagram problem

sort

Map("eilnst", Silent")
Map("eilnst", "Listen")

Reduce("eilnst", [Silent, Listen])

=> Each reducer takes care of each key.

### Remember!

The Map() and Reduce() blocks are the personal loops of the tasks, i.e.,

- 1 Map per partitioned data group.
- 1 Reduce per unique key.

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# Mapper Functions (Reducer Similar)



### **Mapper v2.8 Functions**

- •protected void **setup**(org.apache.hadoop.mapreduce.Mapper.Context context) throws <u>IOException</u>, <u>InterruptedException</u>
  - Called once at the beginning of the task.
- •protected void map(KEYIN key, VALUEIN value,
  org.apache.hadoop.mapreduce.Mapper.Context context) throws IOException,
  InterruptedException
  - Called once for each key/value pair in the input split. Most applications should override this, but the default is the identity function (k,v) => (k,v).
- •protected void **cleanup(**org.apache.hadoop.mapreduce.Mapper.Context context) throws IOException, InterruptedException
  - Called once at the end of the task.
- •public void **run**(org.apache.hadoop.mapreduce.Mapper.Context context) throws <a href="IOException">IOException</a>, <a href="InterruptedException">InterruptedException</a>
  - Expert users can override this method for more complete control over the execution of the Mapper.
- •Methods inherited from class java.lang.Object: clone, equals, finalize, getClass, hashCode, toString, thread control: notify, notifyAll, wait
- •How to run the code More details:

# Word Count Pseudocode (easier for next slides)



"Hello World" in "Map Reduce"

```
1: class MAPPER
                                                                Both MAPPER and
       method MAP(offset a, line I)
2:
                                                           REDUCER could skip EMIT(),
           for all term t \in \text{line } I do
3:
                                                                e.g., the case of filter
4:
               EMIT(term t, count 1)
1: class Reducer
       method REDUCE(term t, counts [c_1, c_2, \ldots])
                                                                  Iterator <> TABLE
2:
                                                               (we need to begin from
3:
           sum \leftarrow 0
                                                                 the beginning in the
           for all count c \in \text{counts } [c_1, c_2, \ldots] do
4:
                                                                 iterator if we want to
5:
               sum \leftarrow sum + c
                                                               rewind, iterator.reset())
           EMIT(term t, count sum)
6:
```

# Example: Mean Computation



#### **Problem:**

- We have a large dataset where input keys are strings and input values are integers (e.g., <a href="http://www.cs.ucy.ac.cy">http://www.cs.ucy.ac.cy</a>, 10s)
- We wish to compute the mean of all integers associated with the same key
- Almost identical to "word count" (now "word average")!

```
1. class MAPPER
       method MAP(string t, integer r)
                                                      Identity emit function (could
           EMIT(string t, integer r)
                                                      be omitted as it is default)
3:
1: class Reducer
       method REDUCE(string t, integers [r_1, r_2, \ldots])
2:
          sum \leftarrow 0
3:
          cnt \leftarrow 0
4:
          for all integer r \in \text{integers } [r_1, r_2, \ldots] do
5:
6:
              sum \leftarrow sum + r
                                             Note!: A good idea here would be to use a
7:
              cnt \leftarrow cnt + 1
                                          symmetric hash (e.g., MD5), i.e., key=md5(URL)
          r_{ava} \leftarrow sum/cnt
8:
                                                instead of key=URL to minimize string
           EMIT(string t, integer r_{ava})
9:
                                                  comparisons. At the end we could
                                                   URL = reverse-md5(md5(URL))
                                                                                          16-9
```

## Operational Issues

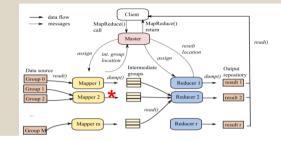


### Aspects that are not under the control of the designer

- Where a mapper or reducer will run
- When a mapper or reducer begins or finishes
- Which input key-value pairs are processed by a specific mapper
- Which intermediate key-value pairs are processed by a specific reducer

### Aspects that can be controlled

- Construct data structures (intermediate results) as keys and values
- Execute user-specified initialization (setup()) and termination code (cleanup()) for mappers and reducers
- Preserve state across multiple input and intermediate keys in mappers and reducers (in-memory combiners – discussed next)
- Control the sort order of intermediate keys, and therefore the order in which a reducer will encounter particular keys.
- Control the partitioning of the key space, and therefore the set of keys that will be encountered by a particular reducer.



# Combiners (optional)



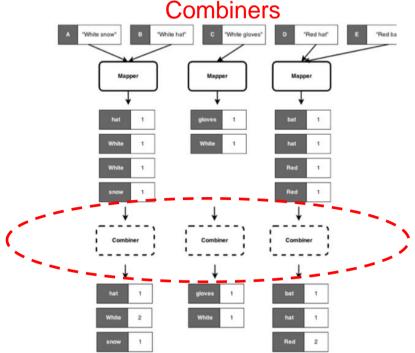
- Combiners are a general mechanism to reduce the amount of intermediate data (after the map task)
  - They could be thought of as "mini-reducers" before data is shipped to reducers
  - Reduce the number and size of key-value pairs to be shuffled
- Back to our running example: word count
  - Combiners aggregate term counts across documents processed by each map task
  - If combiners take advantage of all opportunities for local aggregation we have at most  $m \times V$  intermediate key-value pairs
    - *m*: number of mappers
    - V: number of unique terms in the collection
  - Note: due to Zipfian nature of term distribution not all mappers will see all terms.

Frequency

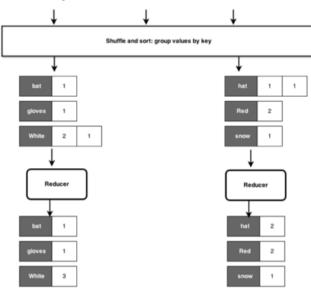
# Combiners (optional)

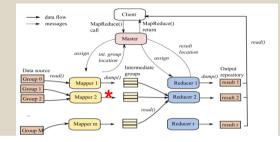


- The use of combiners must be thought carefully
  - In Hadoop, they are optional: the correctness of the algorithm cannot depend on computation (or even execution) of the combiners
  - In Apache Spark, they're mostly automatic



## Shuffling mapreduce.reduce.shuffle.





# Combiners (optional)



- Hadoop does not guarantee combiners to be executed
  - Actually, combiners might only be called if the number of map output records is greater than a threshold, i.e., 4
- Problem: Can we enforce the execution of aggregation at the end of the Map phase?
  - \circ
     \text{Yes, by implementing the aggregation logic in Mapper.}
  - B Not always very good as it the function state:
    - In-memory combining breaks the functional programming paradigm due to state preservation.
    - In-memory combining strictly depends on having sufficient memory to store intermediate results
    - A possible solution: "block" and "flush"
  - Nevertheless, let's see an example...

# In-Memory Combiners (inside Mapper)



```
1: class MAPPER

2: method MAP(offset a, line l)

3: H \leftarrow new HashMap

4: for all term t \in line l do

5: H\{t\} \leftarrow H\{t\} + 1

6: for all term t \in H do

7: EMIT(term t, count H\{t\})
```

# We use a hash map to accumulate intermediate results

- The data structure is also know as "associative array" or "dictionary"
- The array is used to tally up (aggregate) term counts within a single "document"
- The Emit method is called only after all InputRecords have been processed

# Selections (σ) in MapReduce \*\*



- **Revision of** Relational Algebra Operators
  - http://www2.cs.ucy.ac.cy/~dzeina/courses/epl342/schedule.html (Lecture 8 and 9).
- In practice, selections do not need a full-blown **MapReduce implementation** 
  - They can be implemented in the map phase alone
  - Actually, they could also be implemented in the reduce portion!
  - Remember that the input to Reduce is an Iterator (it is constructed) as the packets arrive at the reducer not a fully constructed list on which).
- A MapReduce implementation of  $\sigma_{c}(R)$ 
  - For each tuple t in R, check if t satisfies C
  - If so, emit a key/value pair (t, NULL)
  - (VOID) Identity Reducer

# Projections (π) in MapReduce

### A note on duplicates in projections:

- Relational Algebra ( $\pi$ ) generates NO duplicates (RA operates with sets). Notation we use:  $\pi_{DISTINCT,s}(R)$
- SQL (SELECT): generates duplicates (SQL operates with **multisets).** Notation we use SELECT:  $\pi_s(R)$ . Of course there is also SELECT DISTINCT, again notated with  $\pi_{DISTINCT,s}(R)$

### How to implement Projections in MR?

- $-\pi_s(R)$  (ALL): Keeps Duplicates => Only requires map task.
- $-\pi_{DISTINCTS}(R)$ : Removes Duplicates => Requires map + reduce

### • $\pi_{DISTINCTS}(R)$ Implementation

#### MAP

- For each tuple *t* in *R*, construct a tuple *t*' that contains only the S columns.
- Emit a key/value pair (t', NULL)

#### REDUCE

- Foreach key t' obtained from mappers (t', [<NULL>]), take the key only.

# Unions (U) in MapReduce



R: 1, 2, 2, 2, 3, 4, 4

1, 2, 2, 2, 3, 4, 4, 2, 3, 4, 4, 4, 5

S: 2, 3, 4, 4, 4, 5

R **UNION** S: 1, 2, 3, 4, 5 R **UNION** ALL S<sup>1</sup>

### $R \cup S = \{x \mid x \in R \lor x \in S\}$

Relations R and S must have the same schema!

### A note on duplicates in unions:

- R UALL S: Keeps Duplicates => Requires Map Task only
- R U S: Removes Duplicates => Requires Map + Reduce Task

### Outline of R ∪ S Implementation:

- Map tasks will be assigned chunks from either R or S \*
- Mappers don't do much, just pass by to reducers. Reducers do duplicate elimination (not necessary in R ∪<sub>ALL</sub> S)
- \* Note: Hadoop MapReduce supports reading multiple inputs.

### How to implement R ∪ S in MR?

**MAP:** For each tuple **t** in R and S, emit a key/value pair (**t**, **NULL**) // identity function **REDUCE** 

- Foreach key t' obtained from mappers (t', [<NULL>]), take the key only.
- Emit a key/value pair (t', NULL) // i.e., Either an R tuple or an S tuple.
- Also works when R and/or S have duplicates => Still generates (t', [<NULL>]), 16-17
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# Intersection (n) in MapReduce

### $R \cap S = \{x \mid x \in R \land x \in S\}$

- Relations R and S must have the same schema!
- A note on duplicates in intersections:
  - R ∩ S : Removes Duplicates => Map + Reduce
  - R ∩<sub>ALI</sub> S: Keeps Duplicates=> Map+Reduce (not available in most DBMS)
- Outline of R ∩ S Implementation:
  - Map tasks will be assigned chunks from either R or S
  - Mappers don't do much, just pass by to reducers. Reducers do duplicate elimination (not necessary in R ∩ ALL S)
- How to implement R ∩ S in MR (R, S: no duplicates)

**MAP:** For each tuple **t** in R and S, emit a key/value pair (**t**, **t**)

#### **REDUCE**

- Foreach key t' obtained from the mappers, if (t', [t', t']) (i.e., these 2 entries must have come from R and S) then emit the key/value pair (t', NULL)
- Otherwise, emit nothing // i.e., (t', [t']) OR (t', [<NULL>])
- // If R, S contain duplicates, I must annotate the t in map (t,'R'), (t,'S'), ... to 6-18 EPL@void/self-intersection within R-and within Sarespectively ersity of Cyprus)

Examples:

R: 1, **2**, 2, 2, **3**, **4**, **4** S: **2**, **3**, **4**, **4**, 4, 5

R INTERSECT S: 2, 3, 4
R INTERSECT ALL S: 2, 3, 4, 4

# Set Difference (-) Revision



### $R-S=\{x\mid x\in R \land x\notin S\}$ Note: $R-S\neq S-R$

Relations R and S must have the same schema!

(a)	STL	JDEN1
(a	, 310	DEN

Fn	Ln
Susan *	Yao
Ramesh*	Shah
Johnny	Kohler
Barbara	Jones
Amy	Ford
Jimmy	Wang
Ernest	Gilbert

#### INSTRUCTOR

Fname	Lname
John	Smith
Ricardo	Browne
Susan *	Yao
Francis	Johnson
Ramesh*	Shah

|Instructor|=5

|Student|=7

### |Student-Instructor| = 5

#### (d)

Fn	Ln
Johnny	Kohler
Barbara	Jones
Amy	Ford
Jimmy	Wang
Ernest	Gilbert

### |Instructor – Student| = 3

(e)

Fname	Lname
John	Smith
Ricardo	Browne
Francis	Johnson

16-19

## Set Difference (-) in MR



- Outline of R S Implementation:
  - The map function passes tuples from R and S to the reducer
  - it must inform the reducer whether the tuple came from R or S!

### How to implement R - S in MR?

#### MAP:

For a tuple t in R emit a key/value pair (t, 'R') and for a tuple t in S, emit a key/value pair (t,'S')

#### **REDUCE:**

- For each key t, do the following:
  - If the input is (t,['R']), then emit (t, NULL)
  - If the input is (t,['R', 'S']) or (t,['S', 'R']), or (t,['S']), don't emit anything!

# Join (⋈,⊗) in MapReduce



$$R \otimes_{\text{cattr}} S = \sigma_{\text{cattr}} (R \times S)$$

- This topic is subject to continuous refinements
- There are many JOIN operators and many different implementations
- Let's look at two relations R(A, B) and S(B, C)
- We must find tuples that agree on their B components
- We shall use the B-value of tuples from either relation as the key
- The value will be the other component and the name of the relation
- That way the reducer knows from which relation each tuple is coming from

### How to implement R ⊗ S in MR?

#### MAP:

- For each tuple (a,b) of R emit the key/value pair (b, ('R',a))
- For each tuple (b,c) of S emit the key/value pair (b, ('S',c))

#### **REDUCE:**

- Each key b will be associated to a list of pairs that are either ('R', a) or ('S', c)
- Generate key/value pairs (b,[( $a_1$ ,b, $c_1$ ), ( $a_2$ ,b, $c_2$ ), ..., ( $a_n$ ,b, $c_n$ )] and emit the <u>unique</u> triples (a,b,c) => the final unique step would be best to be implemented with a second MR

job \_see\_assignment! Topics in Databases - Demetris Zeinalipour © (University of Cyprus)

# Aggregation (γ) in MapReduce

- We already discussed Aggregates, remember the Mean Example.
  - Map: The map operation prepares the grouping
  - Reduce: The reducer computes the aggregation.
  - Simplifying assumptions: one grouping attribute and one aggregation function. –
     easy to lift these assumptions and generalize the discussion.
- Different Types of Aggregates :
  - Distributive Aggregates: COUNT,SUM,MAX,MIN,AVG(S/C) => reduce uses a rolling computation of the aggregate.
  - Holistic Aggregates: MEDIAN, MEAN => reduce has to retain all incoming tuples (coming through the iterator) (k,[v1,...,vn]) and then compute the aggregate.
- How to implement γ<sub>A,θ(B)</sub>R in MR?

**MAP:** For a tuple *t* (*a*,*b*,*c*) in *R* emit a key/value pair (*a*,*b*)

**REDUCE:** For each key t, do the following:

- Distributive Aggregates: Apply 0 on the incoming tuples [b1,...bn] on-the-fly
- Holistic Aggregates: Accumulate the incoming tuples in a table. At the end apply on the constructed table.
- Emit the key/value pair (a,x) where  $x = \theta([b1,...,bn])$

# Generalizing MR Operators



- Having developed operators for basic RA operators, one could now develop higher-level declarative languages (e.g., SQL, PIG) that translate into MR jobs!
- Example Apache HUE on top of HIVE (Hadoop / HDFS)

